# Tail Call Elimination in Opensmalltalk

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# Agenda

- What is a Tail Call?
- Tail Call Elimination
- Stack Interpreter Implementation
- Cog VM JIT Implementation
- Results
- Conclusions and Future Work

### • Call followed by a return

#### • Smalltalk Example

- Call to basicNew: is a tail call
- Immediately followed by a return

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# Tail call in Bytecode

• Bytecode for Array class»new:

self
pushTemp: 0
send: basicNew:
returnTop

# Calling new without TCE

#### • Stack after calling new:

new:'s stack frame new:'s argument Sender of new:'s frame

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### • Stack after calling basicNew:

basicNew:'s stack frame basicNew:'s argument new:'s stack frame new:'s argument

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### • Stack after returning from basicNew:

basicNew:'s result new:'s stack frame new:'s argument Sender of new:'s frame

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# Tail Call Elimination

#### • Why return to new:?

• Why keep new:'s stack frame?

# Tail Call Elimination

- Why return to new:?
- Why keep new:'s stack frame?

# Calling new with TCE

new:'s stack frame

new:'s argument

# Calling new with TCE - cont'd

basicNew:'s stack frame

basicNew:'s argument

# Calling new with TCE - cont'd

basicNew:'s return

# Tail Recursion Elimination

### • Special Case of Tail Call Elimination

• Recursive Call is also a Tail Call

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### • Well Known Optimization

- Support in functional languages, CLR, etc.
- Not supported in JVM, Python
- Necessary for functional languages
- Can be useful for OO as well
- In common patterns like Visitor Pattern
- Iteration in Smalltalk

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# Frequency of Tail Calls

### • Static Frequency

Platform Packages	Tail Calls	Total	Percentage
Squeak - All	25162	407971	6.17
Squeak - Compiler	863	8747	9.87

### • Higher Dynamic Frequency

Action	Tail Calls	Total	Percentage
Startup	47669	219054	21.76
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#### • Interpreter only

- Look ahead to next bytecode for return
- Switch to tail call eliminating implementation
- Remove the existing stack frame and arguments
- Pushes the new arguments and calls the next method

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### • Cog VM - levels of inline caching

- No Inline Cache
- Monomorphic Send Sites
- Polymorphic Send Sites
- Megamorphic Send Sites

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- Activate for monomorphic send sites
- Bypass for polymorphic and megamorphic send sites
- Need tail call and non-tail call JIT code for each send
- Copy method lookup code to sender in tail calls

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### Results

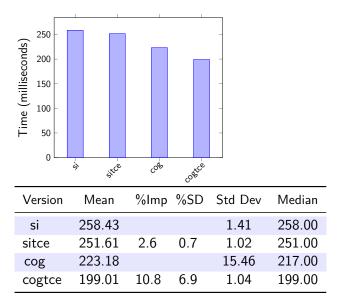
#### • Tail Recursive Tests

• Real World Tests

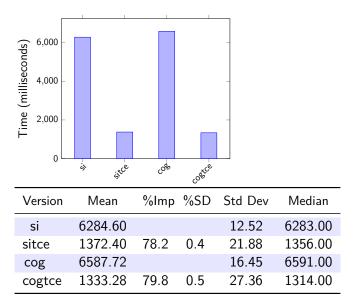
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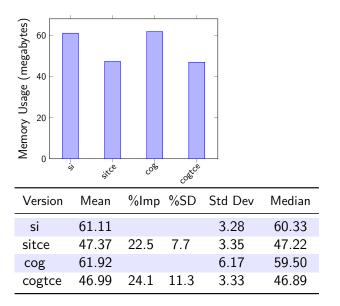
# Factorial $500 \times 1000$



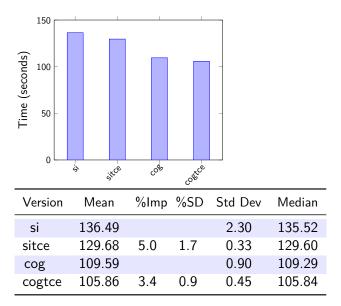
# Factorial 5000



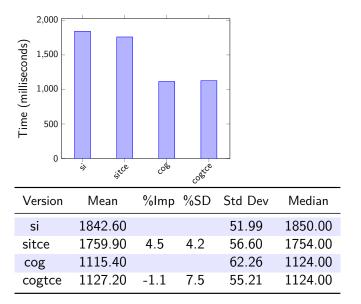
# Factorial 5000 Memory



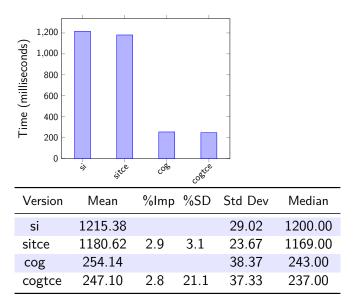
# Compile All Execution



# Browse Number Class Execution



# Method Analyzer Execution



### • Significant improvements in execution time for tail recursive cases

- Improvements in execution time for most general applications
- Memory usage is only significantly affected for deep recursive call chains
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• Reduce redundant code generation for Cog

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# Questions?