

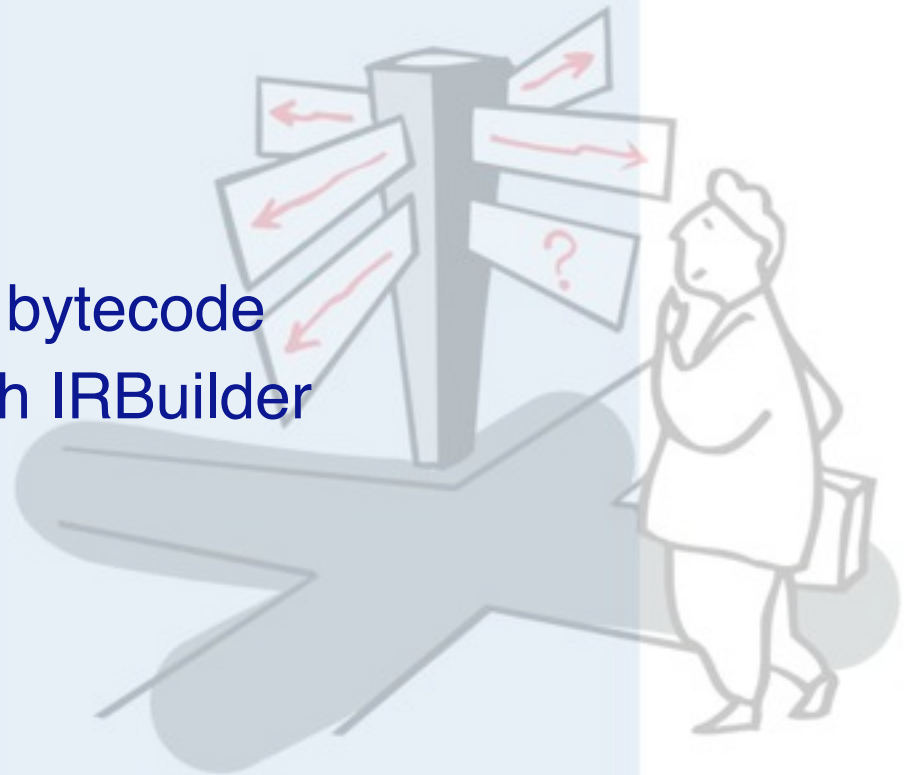
Opal Compiler

Jorge Ressia



Roadmap

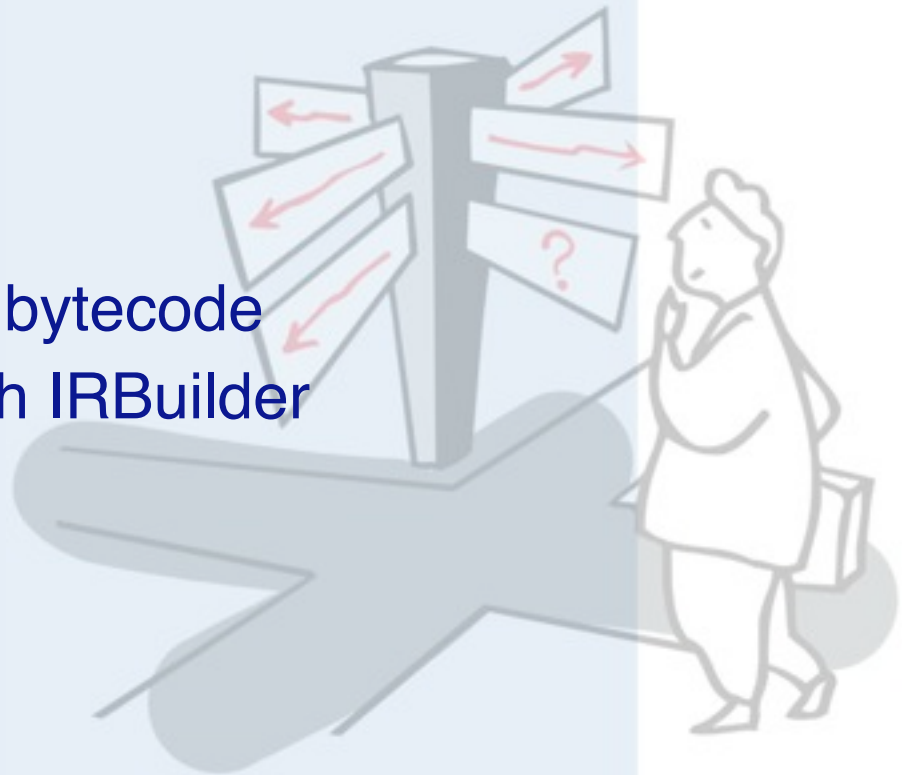
- > The Pharo compiler
- > Introduction to Smalltalk bytecode
- > Generating bytecode with IRBuilder
- > ByteSurgeon



Original material by Marcus Denker

Roadmap

- > **The Pharo compiler**
- > Introduction to Smalltalk bytecode
- > Generating bytecode with IRBuilder
- > ByteSurgeon



The Pharo Compiler

- > Default compiler
 - very old design
 - quite hard to understand
 - hard to modify and extend

What qualities are important in a compiler?

- > Correct code
- > Output runs fast
- > Compiler runs fast
- > Compile time proportional to program size
- > Support for separate compilation
- > Good diagnostics for syntax errors
- > Works well with the debugger
- > Good diagnostics for flow anomalies
- > Consistent, predictable optimization

Why do we care?

- > [ByteSurgeon](#) — Runtime Bytecode Transformation for Smalltalk
- > [ChangeBoxes](#) — Modeling Change as a first-class entity
- > [Reflectivity](#) — Persephone, Geppetto and the rest
- > [Helvetia](#) — Context Specific Languages with Homogeneous Tool Integration
- > [Albedo](#) — A unified approach to reflection.

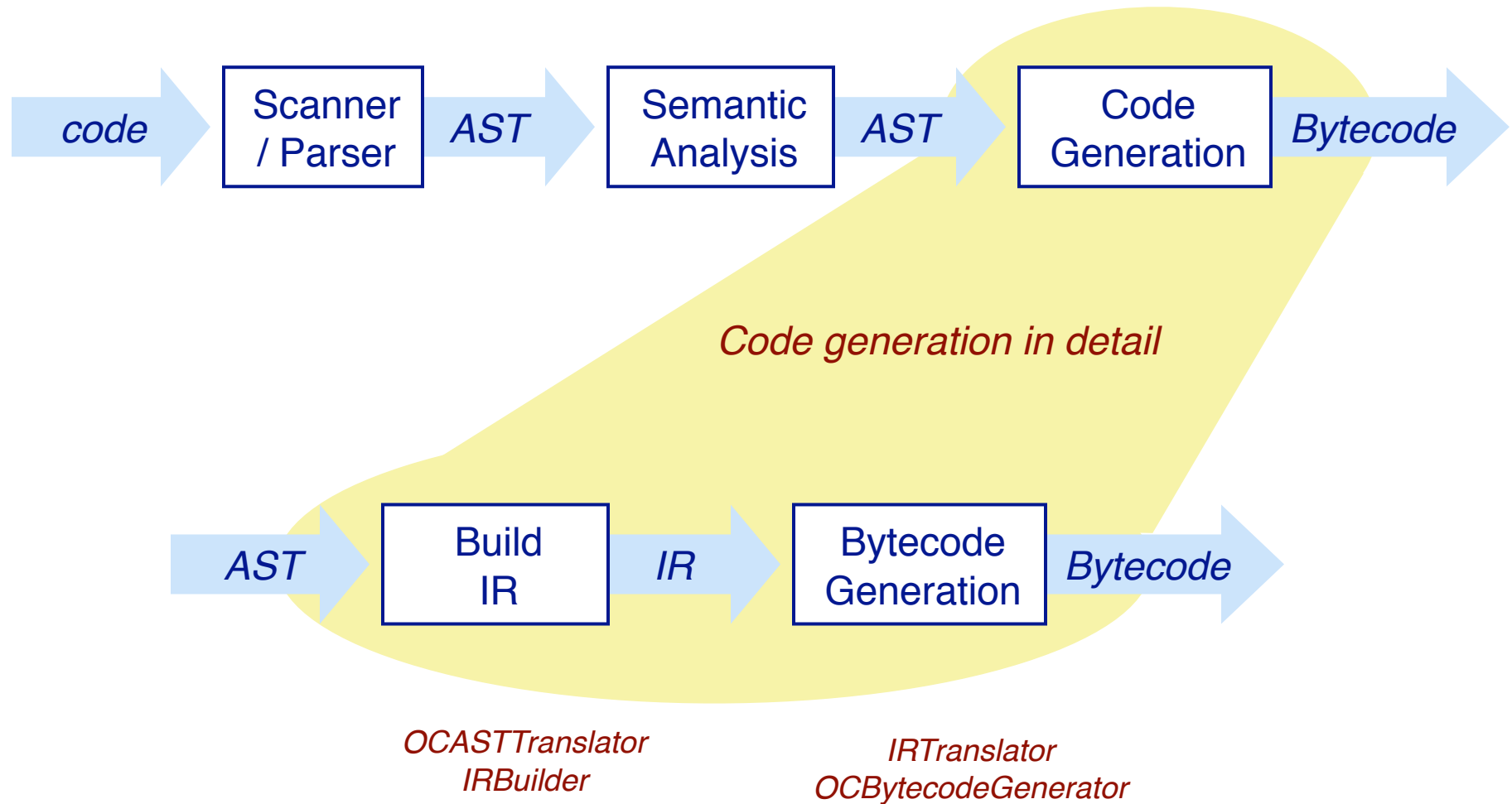
Opal Compiler

- > Opal Compiler for Pharo
 - <http://scg.unibe.ch/research/OpalCompiler>

Opal Compiler

- > Fully reified compilation process:
 - Scanner/Parser (RBParser)
 - *builds AST (from Refactoring Browser)*
 - Semantic Analysis: OCASTSemanticAnalyzer
 - *annotates the AST (e.g., var bindings)*
 - Translation to IR: OCASTTranslator
 - *uses IRBuilder to build IR (Intermediate Representation)*
 - Bytecode generation: IRTranslator
 - *uses OCBytecodeGenerator to emit bytecodes*

Compiler: Overview



Compiler: Design Decisions

- > Every building block of the compiler is implemented as a visitor on the representation.
- > The AST is never changed

Compiler: AST

> AST: Abstract Syntax Tree

- Encodes the Syntax as a Tree
- No semantics yet!
- Uses the RB Tree:
 - *Visitors*
 - *Transformation (replace/add/delete)*
 - *Pattern-directed TreeRewriter*
 - *PrettyPrinter*

```
RBProgramNode
  RBDoItNode
  RBMethodNode
  RBReturnNode
  RBSequenceNode
  RBValueNode
    RBArraryNode
    RBAssignmentNode
    RBBlockNode
    RBCascadeNode
    RBLiteralNode
    RBMessageNode
    RBOptimizedNode
    RBVariableNode
```

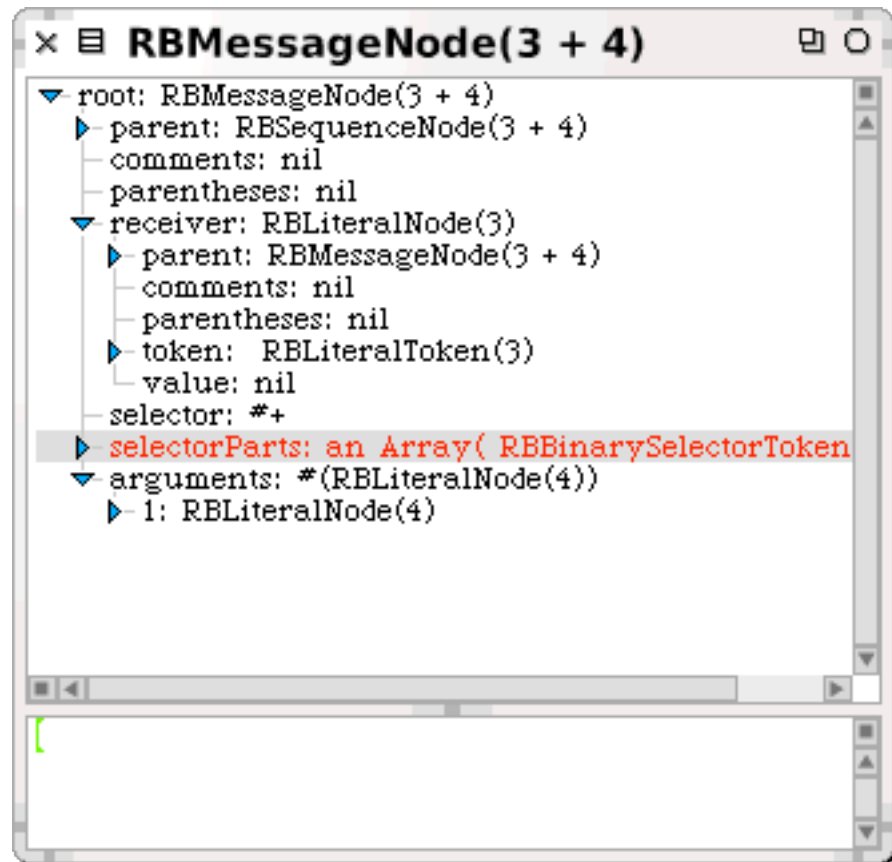
Compiler: Syntax

- > Before: SmaCC: Smalltalk Compiler Compiler
 - Similar to Lex/Yacc
 - SmaCC can build LARL(1) or LR(1) parser
- > Now: RBParser
- > Future: PetitParser

A Simple Tree

RBParser parseExpression: '3+4'

NB: explore it



```
× RBMessageNode(3 + 4)
  ▼ root: RBMessageNode(3 + 4)
    ▶ parent: RBSequenceNode(3 + 4)
      - comments: nil
      - parentheses: nil
    ▼ receiver: RBLiteralNode(3)
      ▶ parent: RBMessageNode(3 + 4)
        - comments: nil
        - parentheses: nil
      ▶ token: RBLiteralToken(3)
        - value: nil
      - selector: #+
      ▶ selectorParts: an Array( RBinarySelectorToken
    ▼ arguments: #(RBLiteralNode(4))
      ▶ 1: RBLiteralNode(4)
```

A Simple Visitor

```
RBProgramNodeVisitor new visitNode: tree
```

Does nothing except
walk through the tree

TestVisitor

```
RBProgramNodeVisitor subclass: #TestVisitor
  instanceVariableNames: 'literals'
  classVariableNames: ''
  poolDictionaries: ''
  category: 'Compiler-AST-Visitors'
```

```
TestVisitor>>acceptLiteralNode: aLiteralNode
  literals add: aLiteralNode value.
```

```
TestVisitor>>initialize
  literals := Set new.
```

```
TestVisitor>>literals
  ^literals
```

```
tree := RBParser parseExpression: '3 + 4'.
(TestVisitor new visitNode: tree) literals
```

```
a Set(3 4)
```

Compiler: Semantics

- > We need to analyze the AST
 - Names need to be linked to the variables according to the scoping rules

- > OCASTSemanticAnalyzer implemented as a Visitor
 - Subclass of RBProgramNodeVisitor
 - Visits the nodes
 - Grows and shrinks scope chain
 - Methods/Blocks are linked with the scope
 - Variable definitions and references are linked with objects describing the variables

Scope Analysis

```
testBlockTemp
| block block1 block2 |
block := [ :arg | [ arg ] ].
block1 := block value: 1.
block2 := block value: 2.
```

Scope Analysis

```
testBlockTemp
| block block1 block2 |
block := [ :arg | [ arg ] ].
block1 := block value: 1.
block2 := block value: 2.
```

```
OCClassScope
  OCInstanceScope
    OCMethodScope 2
      OCBlockScope 3
        OCBlockScope 4
```

Compiler: Semantics

- > OCASTClosureAnalyzer
 - Eliot's Closure analysis: copying vs. tempvector

Closures

```
counterBlock  
  | count |  
  count := 0.  
  ^ [ count := count + 1 ].
```

Closures

- > **Break** the dependency between the block activation and its enclosing contexts for **accessing locals**

Contexts

```
inject: thisValue into: binaryBlock
| nextValue |
nextValue := thisValue.
self
  do: [:each |
    nextValue := binaryBlock
      value: nextValue value: each].
^nextValue
```

Contexts

```
inject: thisValue into: binaryBlock
| indirectTemps |
indirectTemps := Array new: 1.
indirectTemps at: 1 put: thisValue.
" was nextValue := thisValue."
self do:
  [:each |
    indirectTemps
      at: 1
      put: (binaryBlock
        value: (indirectTemps at: 1)
              value: each)].
^indirectTemps at: 1
```

Contexts

```
inject: thisValue into: binaryBlock
| indirectTemps |
indirectTemps := Array new: 1.
indirectTemps at: 1 put: thisValue.
self do: (thisContext
         closureCopy:
             [:each |
              binaryBlockCopy indirectTempsCopy |
              indirectTempsCopy
                at: 1
                put: (binaryBlockCopy
                     value: (indirectTempsCopy at: 1)
                     value: each)]
         copiedValues:
             (Array with: binaryBlock with: indirectTemps)).
^indirectTemps at: 1
```


Closures Analysis

```
| a |  
a := 1.  
[ a ]
```

Closures Analysis

```
| a |  
a := 1.  
[ a ]
```

a is copied

Closures Analysis

```
| index block collection |  
index := 0.  
block := [  
    collection add: [ index ].  
    index := index + 1 ].  
[ index < 5 ] whileTrue: block.
```

Closures Analysis

```
| index block collection |  
index := 0.  
block := [  
    collection add: [ index ].  
    index := index + 1 ].  
[ index < 5 ] whileTrue: block.
```

index is remote

Compiler: Intermediate Representation

- > IR: Intermediate Representation
 - Semantic like Bytecode, but more abstract
 - Independent of the bytecode set
 - IR is a tree
 - IR nodes allow easy transformation
 - Decompilation to RB AST

- > IR is built from AST using OCASTTranslator:
 - AST Visitor
 - Uses IRBuilder

Compiler: Intermediate Representation

```
IRBuilder new
  pushLiteral: 34;
  storeInstVar: 2;
  popTop;
  pushInstVar: 2;
  returnTop;
  ir.
```

```
17 <20> pushConstant: 34
18 <61> popIntoRcvr: 1
19 <01> pushRcvr: 1
20 <7C> returnTop
```

Compiler: Bytecode Generation

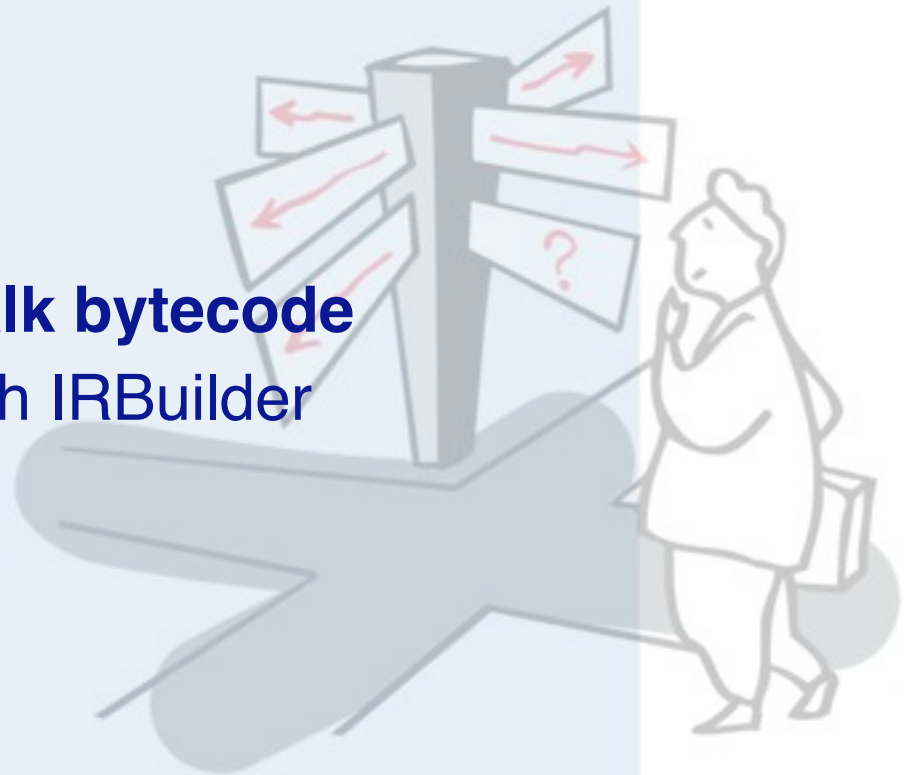
- > IR needs to be converted to Bytecode
 - IRTranslator: Visitor for IR tree
 - Uses OCBytecodeGenerator to generate Bytecode
 - Builds a compiledMethod
 - Details to follow next section

```
testReturn1
| iRMethod aCompiledMethod |
iRMethod := IRBuilder new
  pushLiteral: 1;
  returnTop;
  ir.
```

```
aCompiledMethod := iRMethod compiledMethod.
self should:
  [(aCompiledMethod
    valueWithReceiver: nil
    arguments: #() ) = 1].
```

Roadmap

- > The Pharo compiler
- > **Introduction to Smalltalk bytecode**
- > Generating bytecode with IRBuilder
- > ByteSurgeon



Reasons for working with Bytecode

- > Generating Bytecode
 - Implementing compilers for other languages
 - Experimentation with new language features

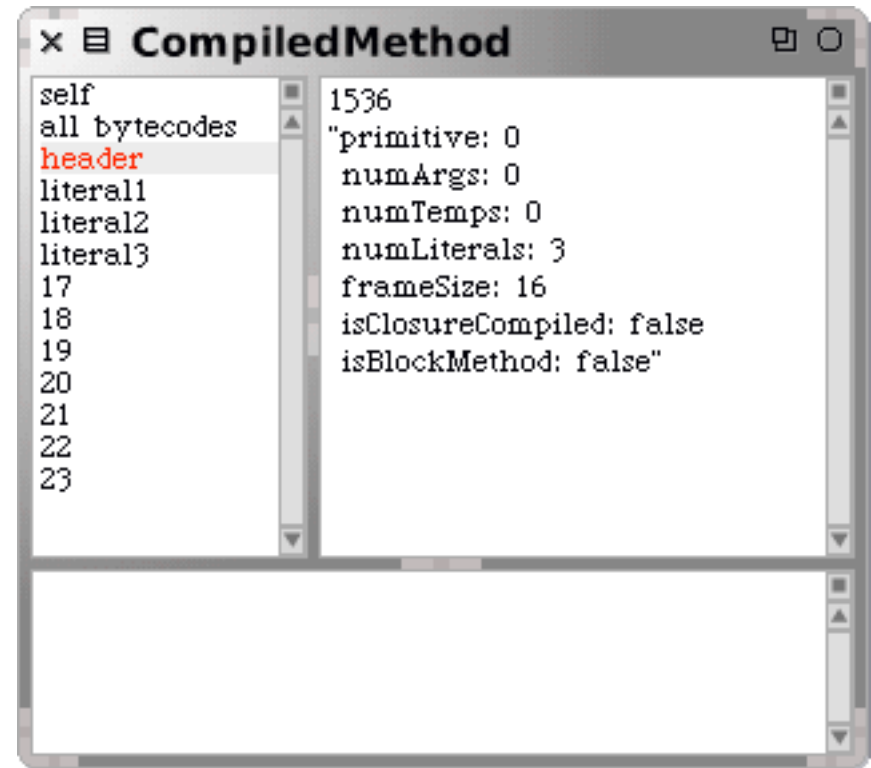
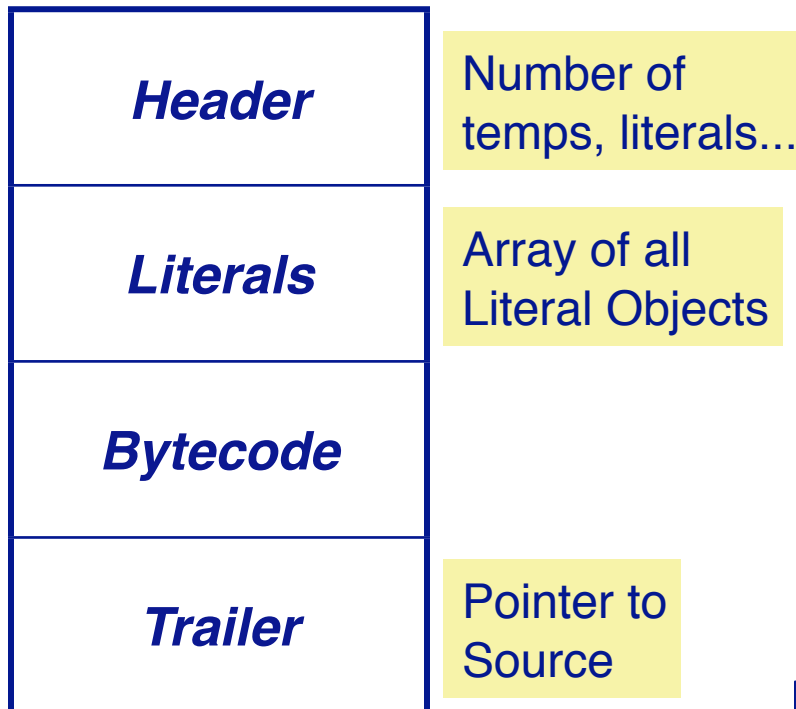
- > Parsing and Interpretation:
 - Analysis (e.g., `self` and `super` sends)
 - Decompilation (for systems without source)
 - Printing of bytecode
 - Interpretation: Debugger, Profiler

The Pharo Virtual Machine

- > Virtual machine provides a virtual processor
 - Bytecode: The “machine-code” of the virtual machine
- > Smalltalk (like Java): Stack machine
 - easy to implement interpreters for different processors
 - most hardware processors are register machines
- > Squeak VM: Implemented in *Slang*
 - Slang: Subset of Smalltalk. (“C with Smalltalk Syntax”)
 - Translated to C

Bytecode in the CompiledMethod

> CompiledMethod format:



```
(Number>>#asInteger) inspect
```

```
(Number methodDict at: #asInteger) inspect
```

Bytecodes: Single or multibyte

> Different forms of bytecodes:

- Single bytecodes:

- *Example: 120: push self*

- Groups of similar bytecodes

- *16: push temp 1*
 - *17: push temp 2*
 - *up to 31*

- Multibyte bytecodes

- *Problem: 4 bit offset may be too small*
 - *Solution: Use the following byte as offset*
 - *Example: Jumps need to encode large jump offsets*

<i>Type</i>	<i>Offset</i>
-------------	---------------

4 bits

4 bits

Example: Number>>asInteger

> Smalltalk code:

```
Number>>asInteger  
    "Answer an Integer nearest  
    the receiver toward zero."  
  
    ^self truncated
```

> Symbolic Bytecode

```
9 <70> self  
10 <D0> send: truncated  
11 <7C> returnTop
```

Example: Step by Step

- > 9 <70> self
 - The receiver (self) is pushed on the stack
- > 10 <D0> send: truncated
 - Bytecode 208: send literal selector 1
 - Get the selector from the first literal
 - start message lookup in the class of the object that is on top of the stack
 - result is pushed on the stack
- > 11 <7C> returnTop
 - return the object on top of the stack to the calling method

Pharo Bytecode

- > 256 Bytecodes, four groups:
 - Stack Bytecodes
 - *Stack manipulation: push / pop / dup*
 - Send Bytecodes
 - *Invoke Methods*
 - Return Bytecodes
 - *Return to caller*
 - Jump Bytecodes
 - *Control flow inside a method*

Stack Bytecodes

- > Push values on the stack
 - e.g., temps, instVars, literals
 - e.g: 16 - 31: push instance variable
- > Push Constants
 - False/True/Nil/1/0/2/-1
- > Push `self`, `thisContext`
- > Duplicate top of stack
- > Pop

Sends and Returns

- > Sends: receiver is on top of stack
 - Normal send
 - Super Sends
 - Hard-coded sends for efficiency, e.g. +, -

- > Returns
 - Return top of stack to the sender
 - Return from a block
 - Special bytecodes for return `self`, `nil`, `true`, `false` (for efficiency)

Jump Bytecodes

- > Control Flow inside one method
 - Used to implement control-flow efficiently
 - Example:

```
^ 1<2 ifTrue: ['true']
```

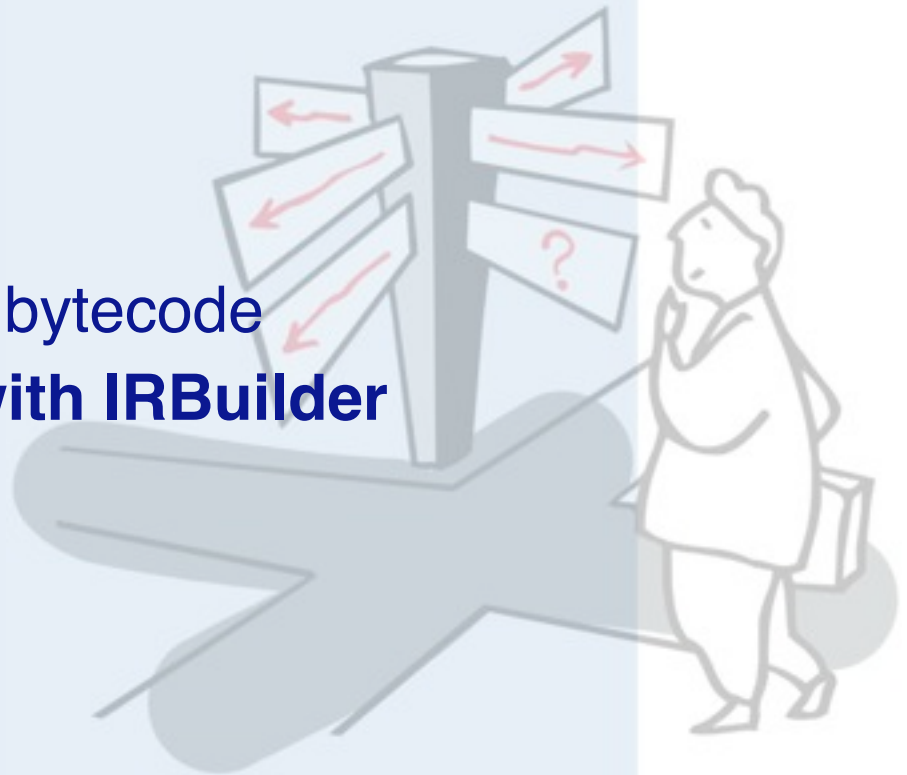
```
9 <76> pushConstant: 1
10 <77> pushConstant: 2
11 <B2> send: <
12 <99> jumpFalse: 15
13 <20> pushConstant: 'true'
14 <90> jumpTo: 16
15 <73> pushConstant: nil
16 <7C> returnTop
```

Closure Bytecode

- > 138 Push (Array new: k)/Pop k into: (Array new: j)
- > 140 Push Temp At k In Temp Vector At: j
- > 141 Store Temp At k In Temp Vector At: j
- > 142 Pop and Store Temp At k In Temp Vector At: j
- > 143 Push Closure Num Copied I Num Args k BlockSize j

Roadmap

- > The Pharo compiler
- > Introduction to Smalltalk bytecode
- > **Generating bytecode with IRBuilder**
- > ByteSurgeon



Generating Bytecode

- > IRBuilder: A tool for generating bytecode
 - Part of the OpalCompiler

- > Like an Assembler for Pharo

IRBuilder: Simple Example

> *Number*>>asInteger

```
iRMethod := IRBuilder new
  pushReceiver;    "push self"
  send: #truncated;
  returnTop;
  ir.

aCompiledMethod := iRMethod compiledMethod.

aCompiledMethod valueWithReceiver:3.5
  arguments: #()
```

3

IRBuilder: Stack Manipulation

- > popTop
 - remove the top of stack
- > pushDup
 - push top of stack on the stack
- > pushLiteral:
- > pushReceiver
 - push self
- > pushThisContext

IRBuilder: Symbolic Jumps

> Jump targets are resolved:

> Example: `false ifTrue: ['true'] ifFalse: ['false']`

```
iRMethod := IRBuilder new
  pushLiteral: false;
  jumpAheadTo: #false if: false;
  pushLiteral: 'true';           "ifTrue: ['true']"
  jumpAheadTo: #end;
  jumpAheadTarget: #false;
  pushLiteral: 'false';         "ifFalse: ['false']"
  jumpAheadTarget: #end;
  returnTop;
  ir.
```


IRBuilder: Instance Variables

- > Access by offset
- > Read: pushInstVar:
 - receiver on top of stack
- > Write: storeInstVar:
 - value on stack
- > Example: set the first instance variable to 2

```
iRMethod := IRBuilder new
  pushLiteral: 2;
  storeInstVar: 1;
  pushReceiver;           "self"
  returnTop;
  ir.

aCompiledMethod := iRMethod compiledMethod.
aCompiledMethod valueWithReceiver: 1@2 arguments: #()
```

IRBuilder: Temporary Variables

- > Accessed by name
- > Define with addTemp: / addTemps:
- > Read with pushTemp:
- > Write with storeTemp:
- > Example:
 - set variables a and b, return value of a

```
iRMethod := IRBuilder new
  addTemps: #(a b);
  pushLiteral: 1;
  storeTemp: #a;
  pushLiteral: 2;
  storeTemp: #b;
  pushTemp: #a;
  returnTop;
  ir.
```

IRBuilder: Sends

> normal send

```
builder pushLiteral: 'hello'  
builder send: #size;
```

> super send

```
...  
builder send: #selector toSuperOf: aClass;
```

- The second parameter specifies the class where the lookup starts.

IRBuilder: Example

```
OCInstanceVar>>emitStore: methodBuilder  
methodBuilder storeInstVar: index
```

IRBuilder: Example

```
OCInstanceVar>>emitStore: methodBuilder  
methodBuilder  
    pushReceiver;  
    pushLiteral: index;  
    send: #instVarAt
```

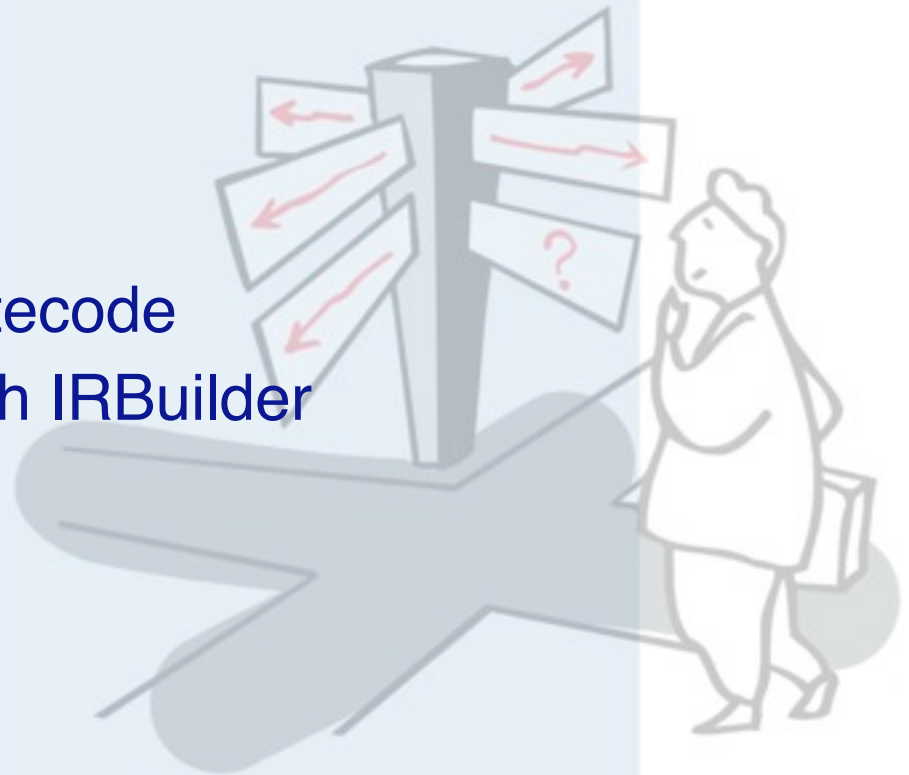
IRBuilder: Example

```
OCInstanceVar>>emitStore: methodBuilder  
methodBuilder  
    pushReceiver;  
    pushLiteral: index;  
    send: #instVarAt:
```

This is global and we do not have much control

Roadmap

- > The Pharo compiler
- > Introduction to Pharo bytecode
- > Generating bytecode with IRBuilder
- > **ByteSurgeon**



ByteSurgeon

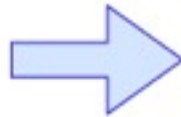
- > Library for bytecode transformation in Smalltalk
- > Full flexibility of Smalltalk Runtime
- > Provides high-level API
- > For Pharo, but portable

- > Runtime transformation needed for
 - Adaptation of running systems
 - Tracing / debugging
 - New language features (MOP, AOP)

Example: Logging

- > Goal: logging message send.
- > First way: Just edit the text:

```
example  
  self test.
```



```
example  
  Transcript show: 'sending #test'.  
  self test.
```

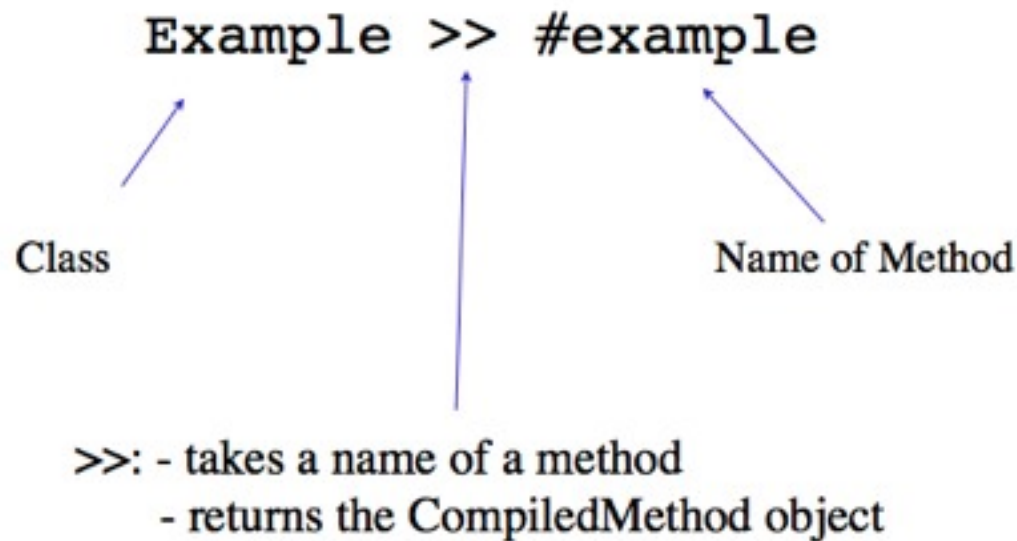
Logging with ByteSurgeon

- > Goal: Change the method without changing program text
- > Example:

```
(Example>>#example)instrumentSend: [:send |  
  send insertBefore:  
    'Transcript show: ''sending #test'' '  
]
```

Logging: Step by Step

```
(Example>>#example)instrumentSend: [:send |
  send insertBefore:
    'Transcript show: ''sending #test'' '.
]
```



Logging: Step by Step

```
(Example>>#example)instrumentSend: [:send |  
  send insertBefore:  
    'Transcript show: ''sending #test'' '.  
]
```

- > instrumentSend:
 - takes a block as an argument
 - evaluates it for all send bytecodes

Logging: Step by Step

```
(Example>>#example)instrumentSend: [:send |  
  send insertBefore:  
    'Transcript show: ''sending #test'' '  
]
```

- > The block has one parameter: send
- > It is executed for each send bytecode in the method

Logging: Step by Step

```
(Example>>#example)instrumentSend: [:send |  
  send insertBefore:  
    'Transcript show: ''sending #test'' '.  
]
```

- > Objects describing bytecode understand how to insert code
 - insertBefore
 - insertAfter
 - replace

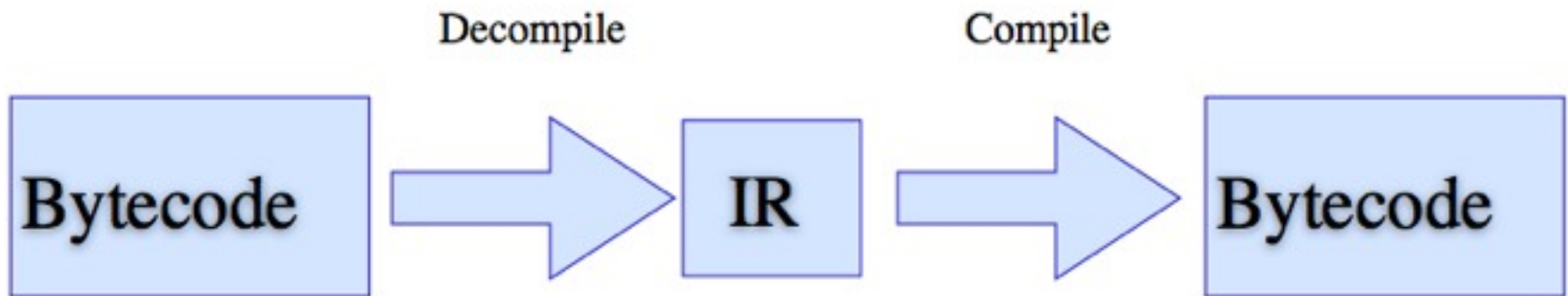
Logging: Step by Step

```
(Example>>#example)instrumentSend: [:send |  
  send insertBefore:  
    'Transcript show: ''sending #test'' '  
]  
]
```

- > The code to be inserted.
- > Double quoting for string inside string
 - *Transcript show: 'sending #test'*

Inside ByteSurgeon

- > Uses IRBuilder internally



- > Transformation (Code inlining) done on IR

ByteSurgeon Usage

> On Methods or Classes:

```
MyClass instrument: [..... ].  
(MyClass>>#myMethod) instrument: [..... ].
```

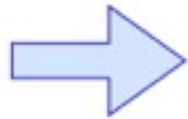
> Different instrument methods:

- instrument:
- instrumentSend:
- instrumentTempVarRead:
- instrumentTempVarStore:
- instrumentTempVarAccess:
- same for InstVar

Advanced ByteSurgeon

- > Goal: extend a send with after logging

```
example  
  self test.
```



```
example  
  self test.  
  Logger logSendTo: self.
```

Advanced ByteSurgeon

- > With ByteSurgeon, something like:

```
(Example>>#example)instrumentSend: [:send |  
  send insertAfter:  
    'Logger logSendTo: ?' .  
]
```

- > How can we access the receiver of the send?
- > Solution: Metavariable

Advanced ByteSurgeon

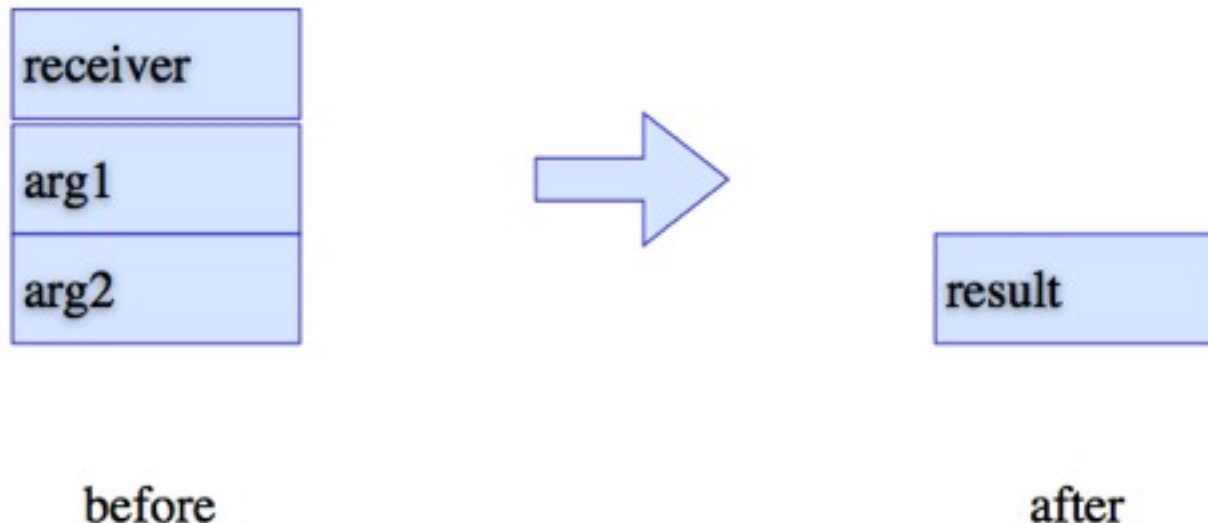
- > With Bytesurgeon, something like:

```
(Example>>#example)instrumentSend: [:send |  
  send insertAfter:  
    'Logger logSendTo: <meta: #receiver>' .  
]
```

- > How can we access the receiver of the send?
- > Solution: Metavariable

Implementation Metavariables

- > Stack during send:



- > Problem I: After send, receiver is not available
- > Problem II: Before send, receiver is deep in the stack

Implementation Metavariables

- > Solution: ByteSurgeon generates preamble
 - Pop the arguments into temps
 - Pop the receiver into temps
 - Rebuild the stack
 - Do the send
 - Now we can access the receiver even after the send

Implementation Metavariables

25 <70> self

26 <81 40> storeIntoTemp: 0

28 <D0> send: test

29 <41> pushLit: Transcript

30 <10> pushTemp: 0

31 <E2> send: show:

32 <87> pop

33 <87> pop

34 <78> returnSelf

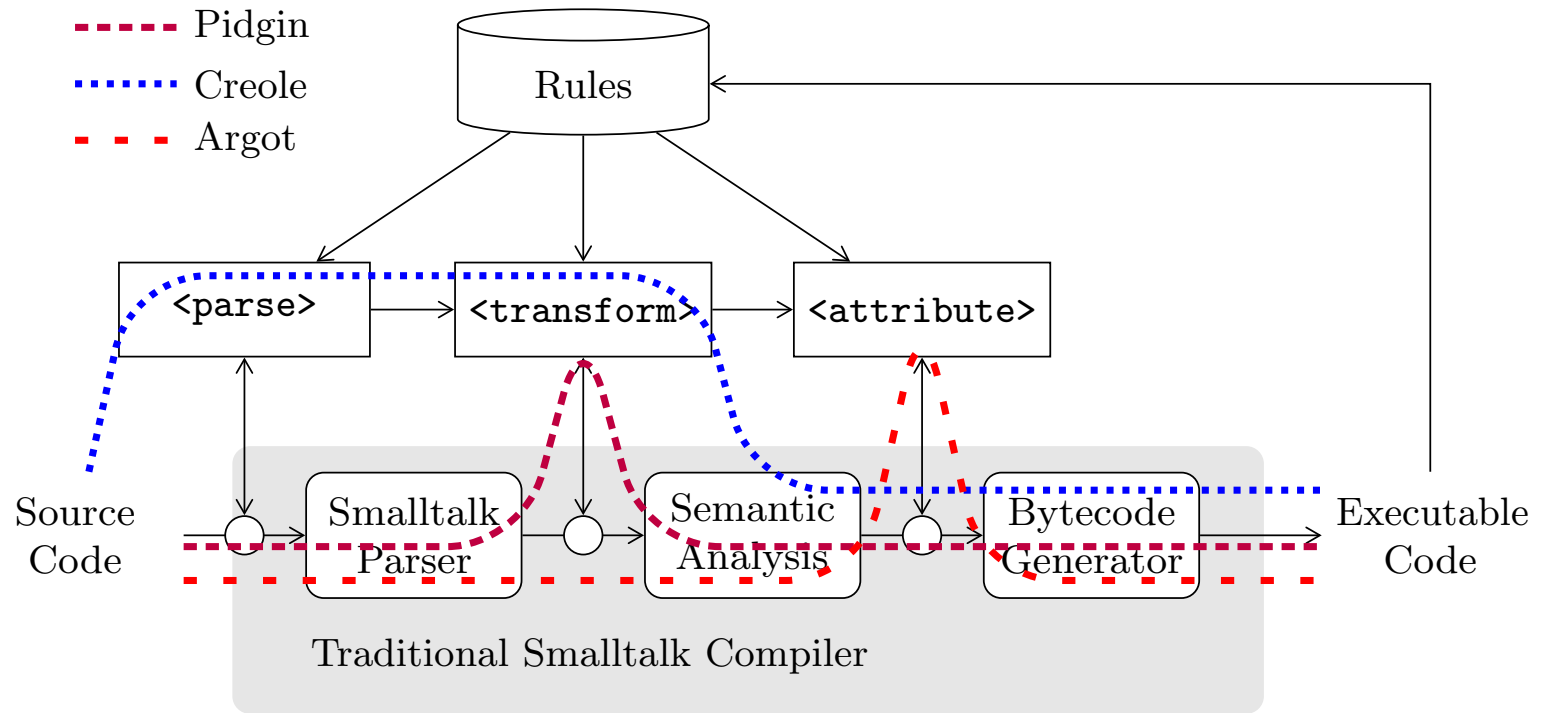
Preamble

Inlined Code

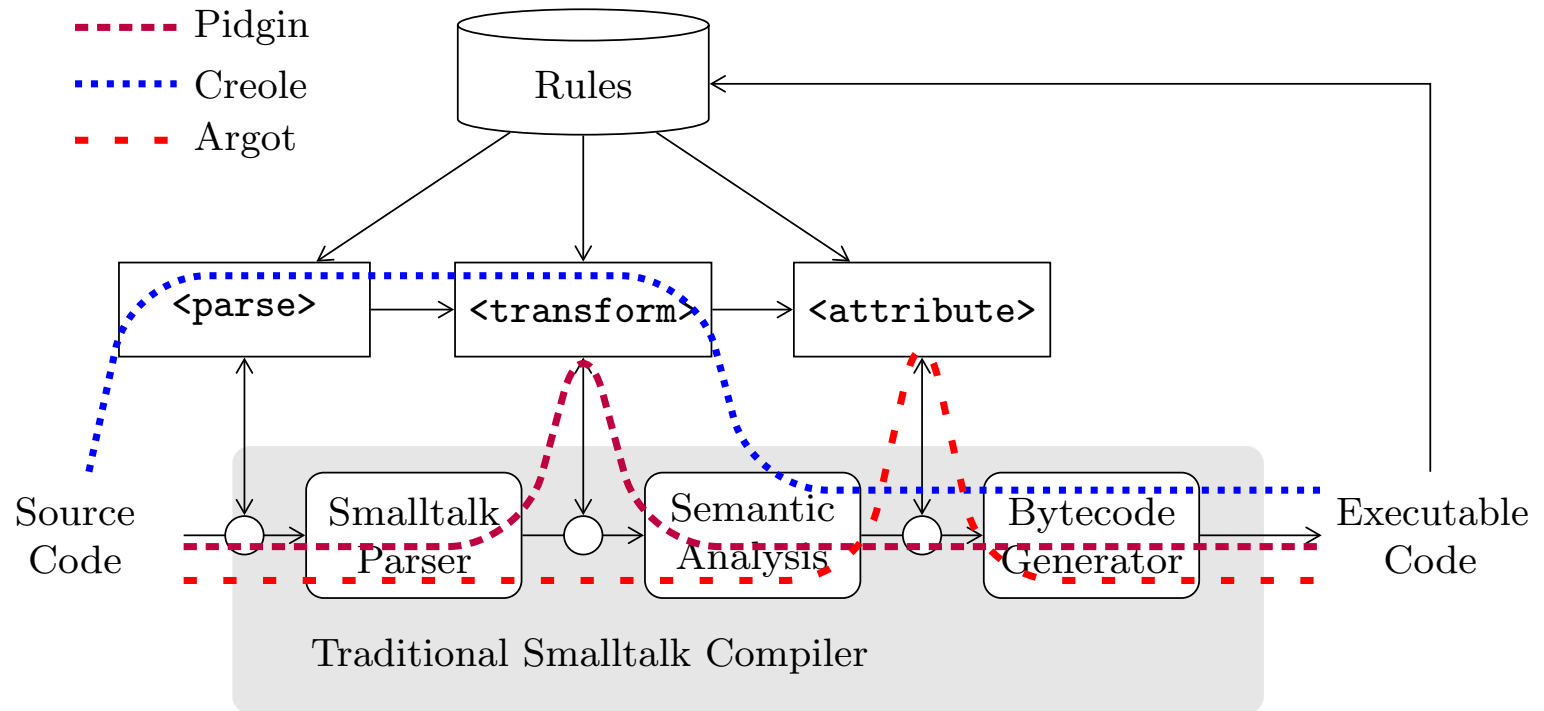
Why do we care?

- > Helvetia — Context Specific Languages with Homogeneous Tool Integration
- > Reflectivity — Unanticipated partial behavioral reflection.
- > Albedo — A unified approach to reflection.

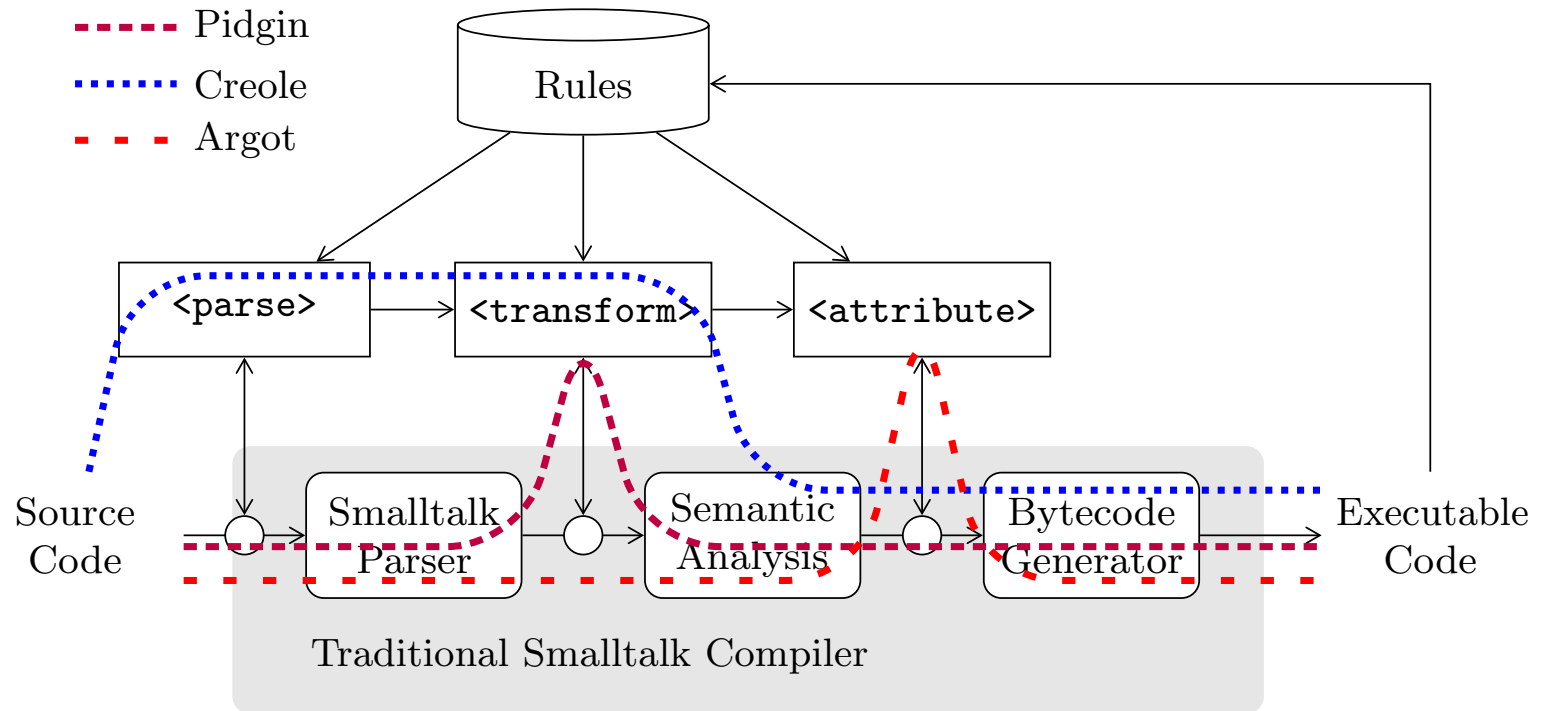
Helvetia



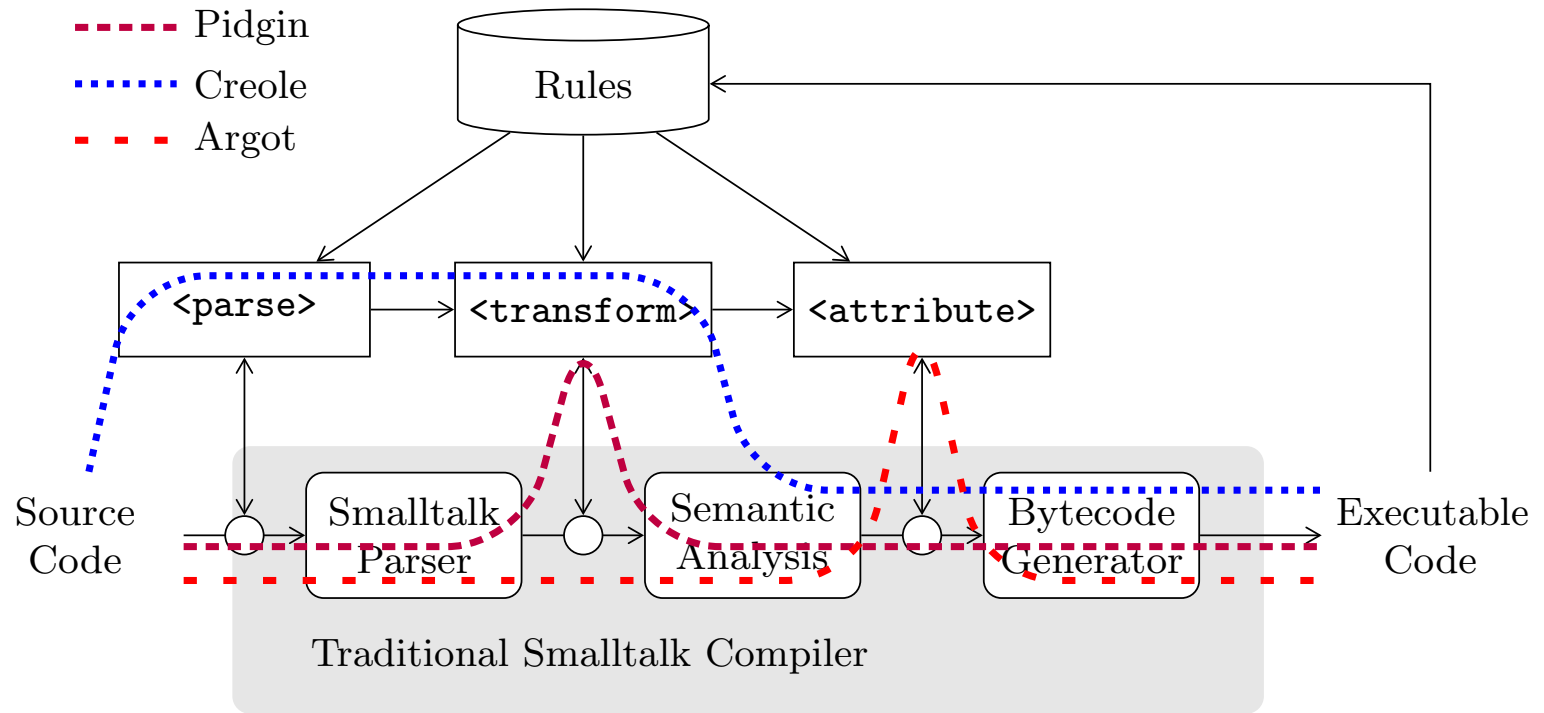
Helvetia



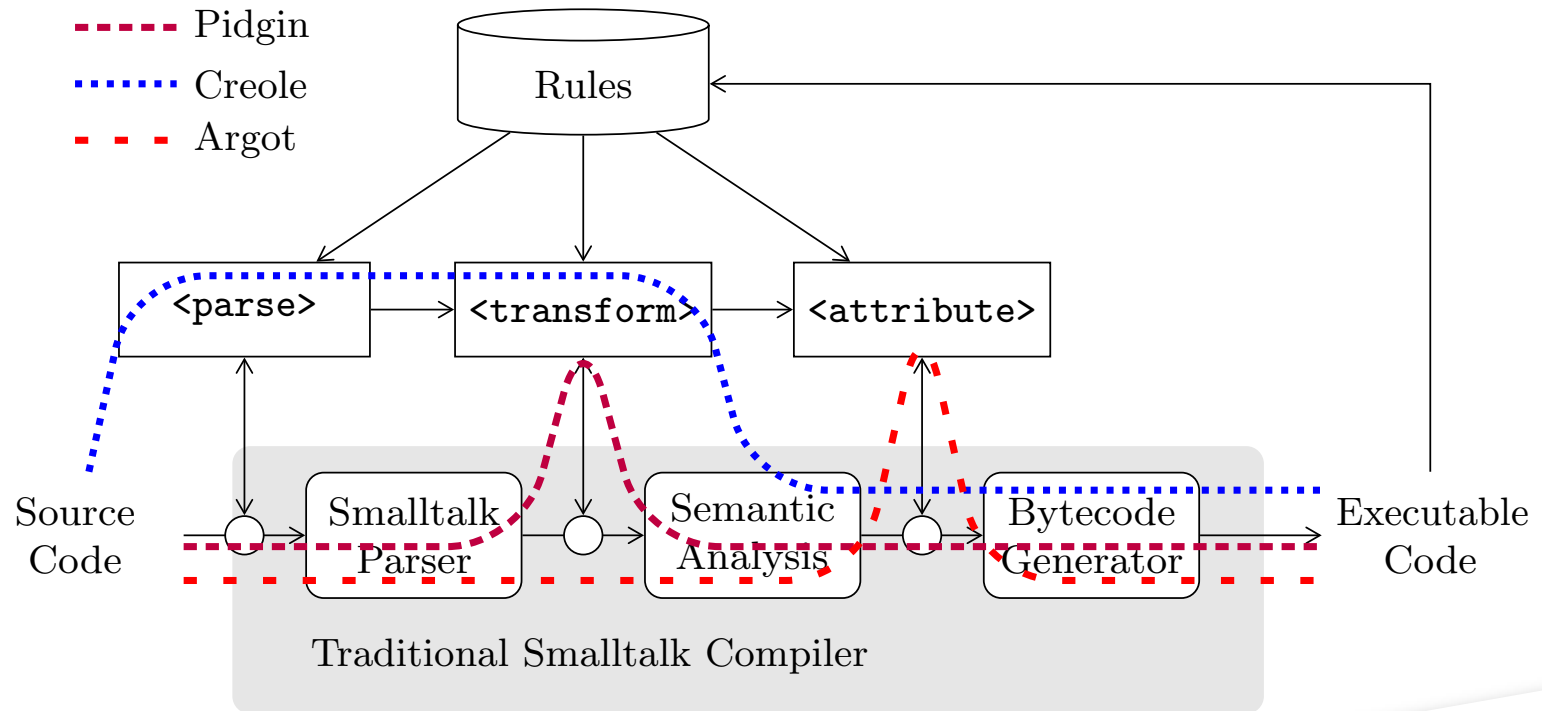
Helvetia



Helvetia

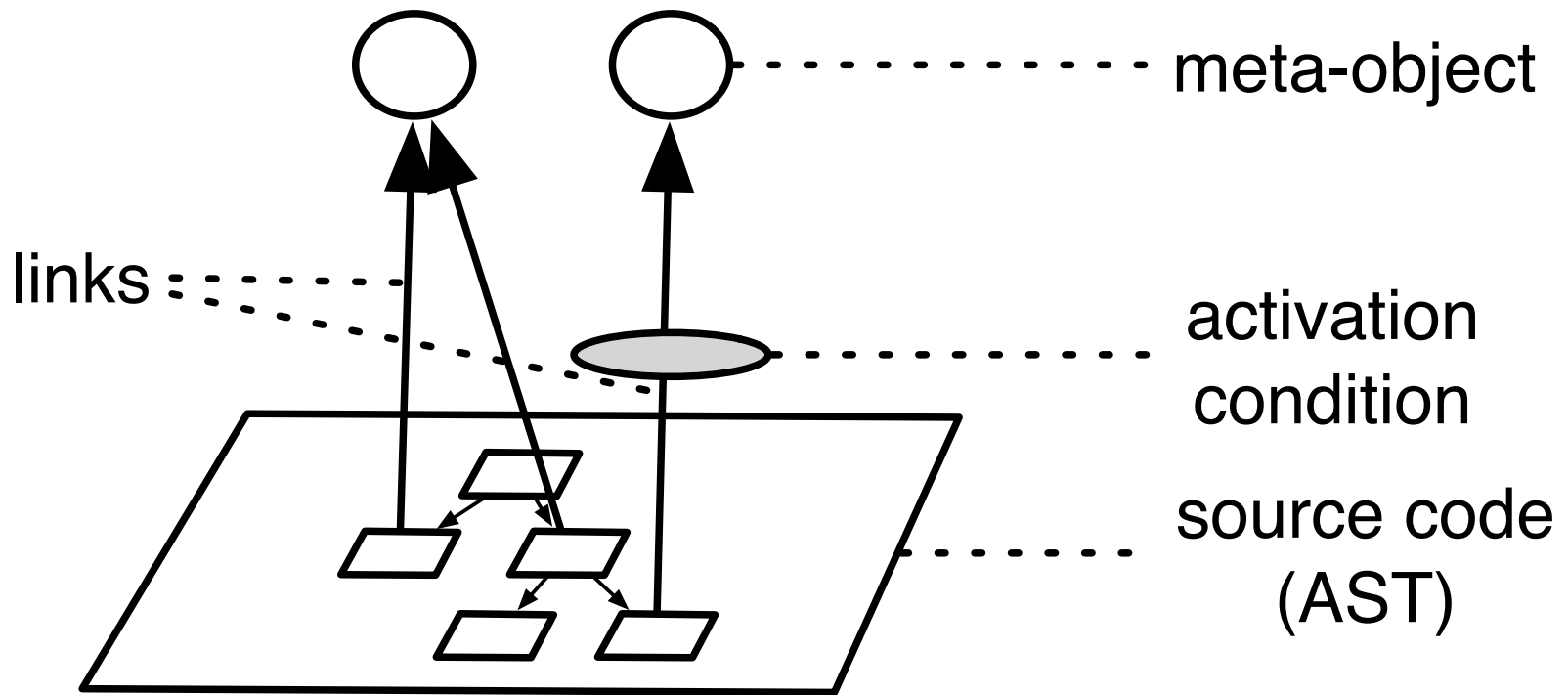


Helvetia

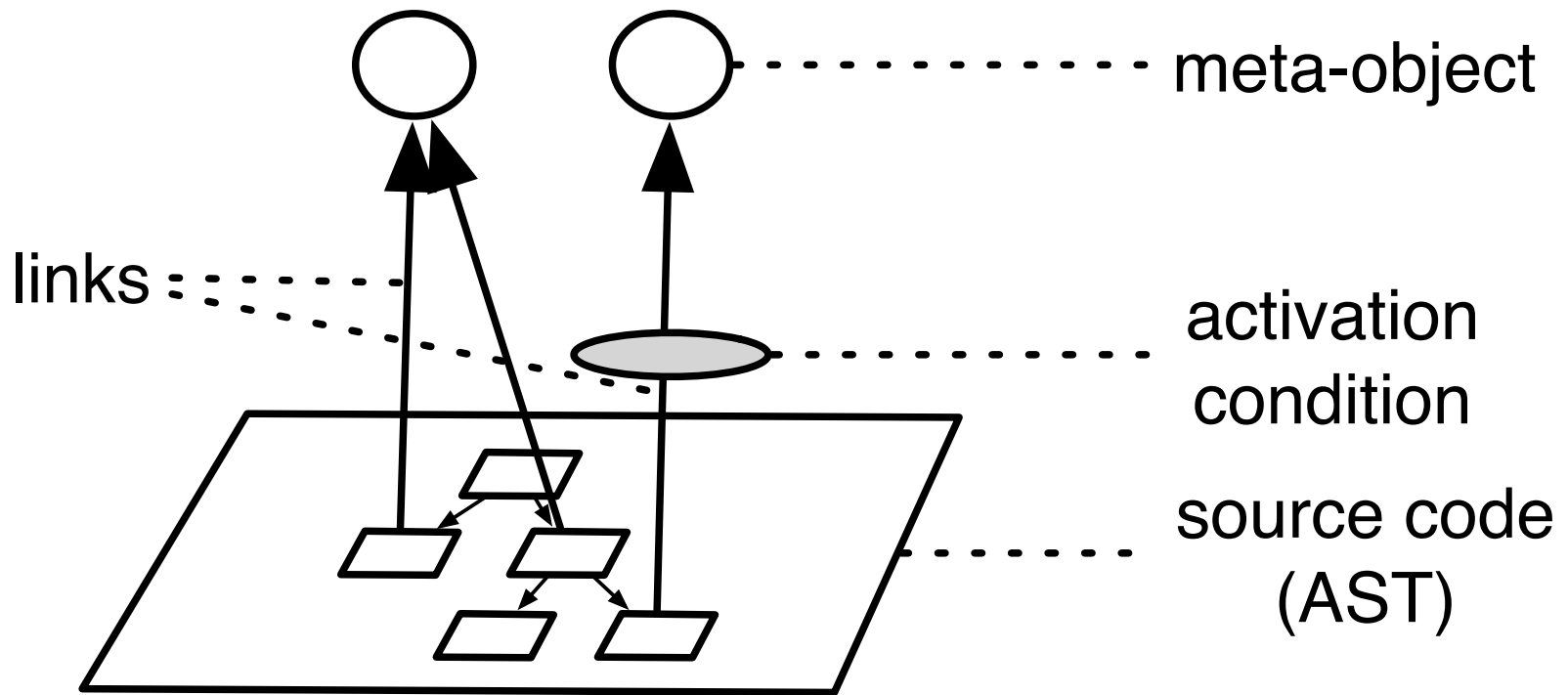


Helvetia
Renggli 2010

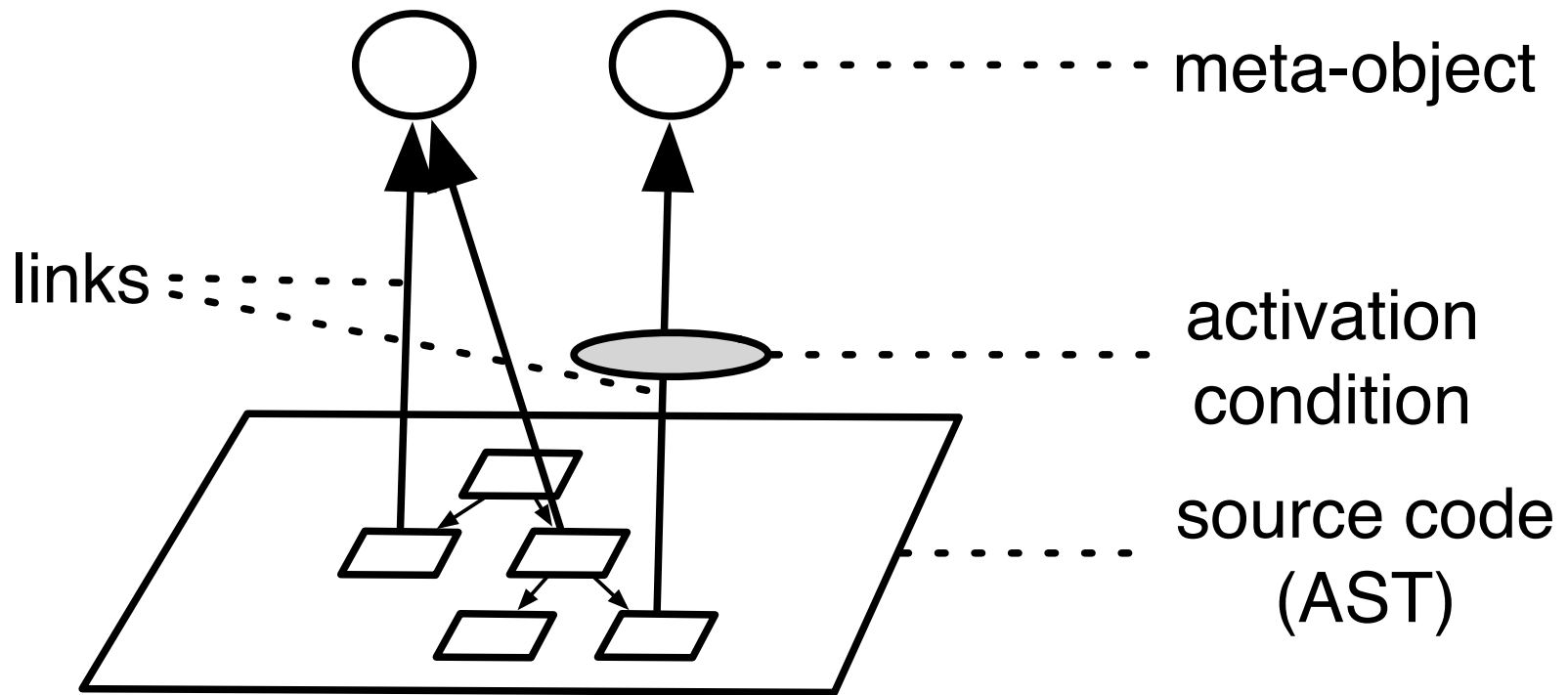
Reflectivity



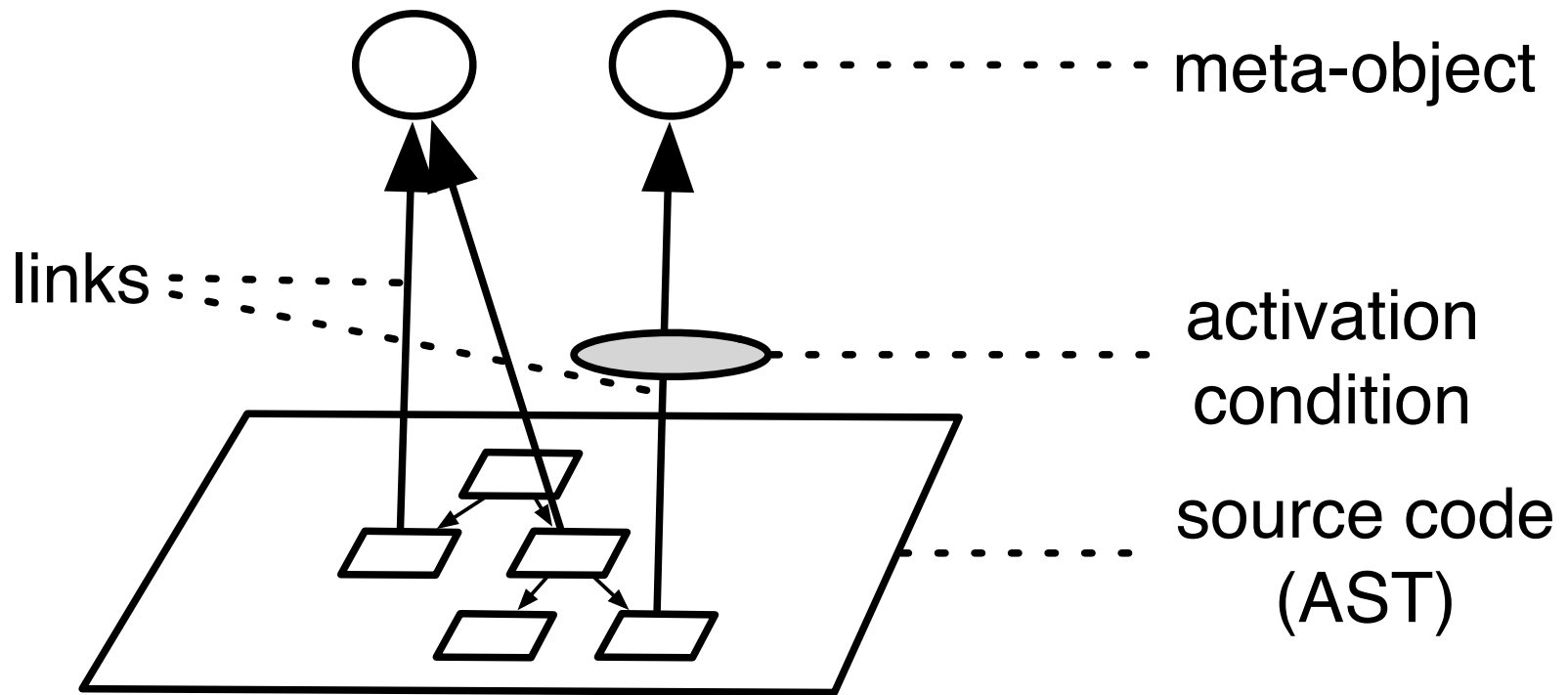
Reflectivity



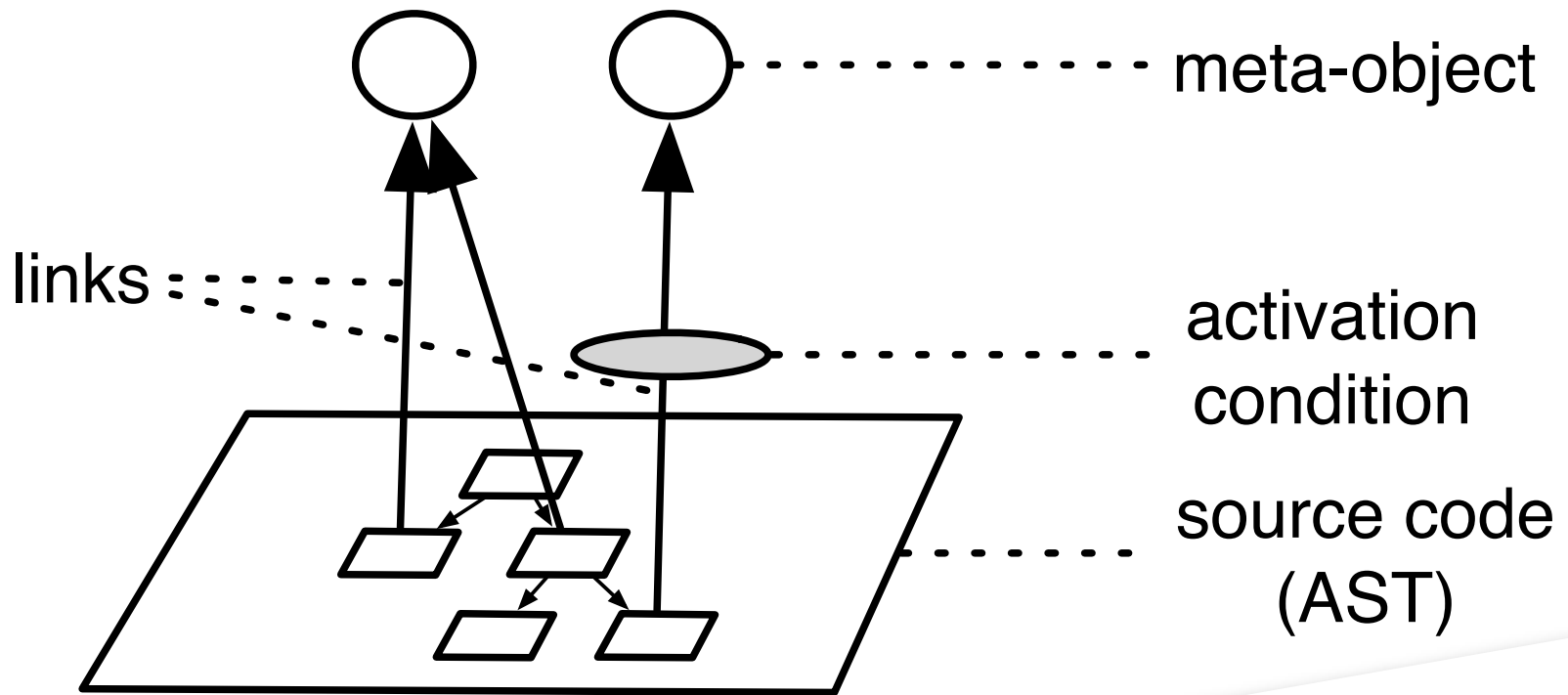
Reflectivity



Reflectivity



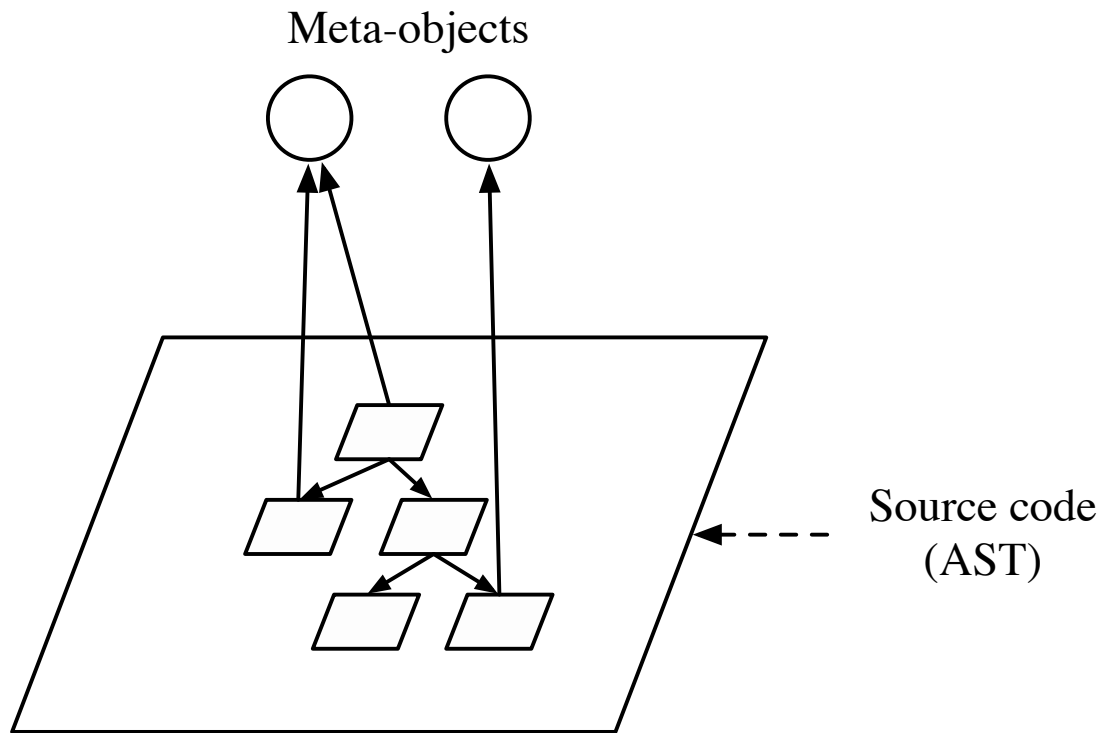
Reflectivity



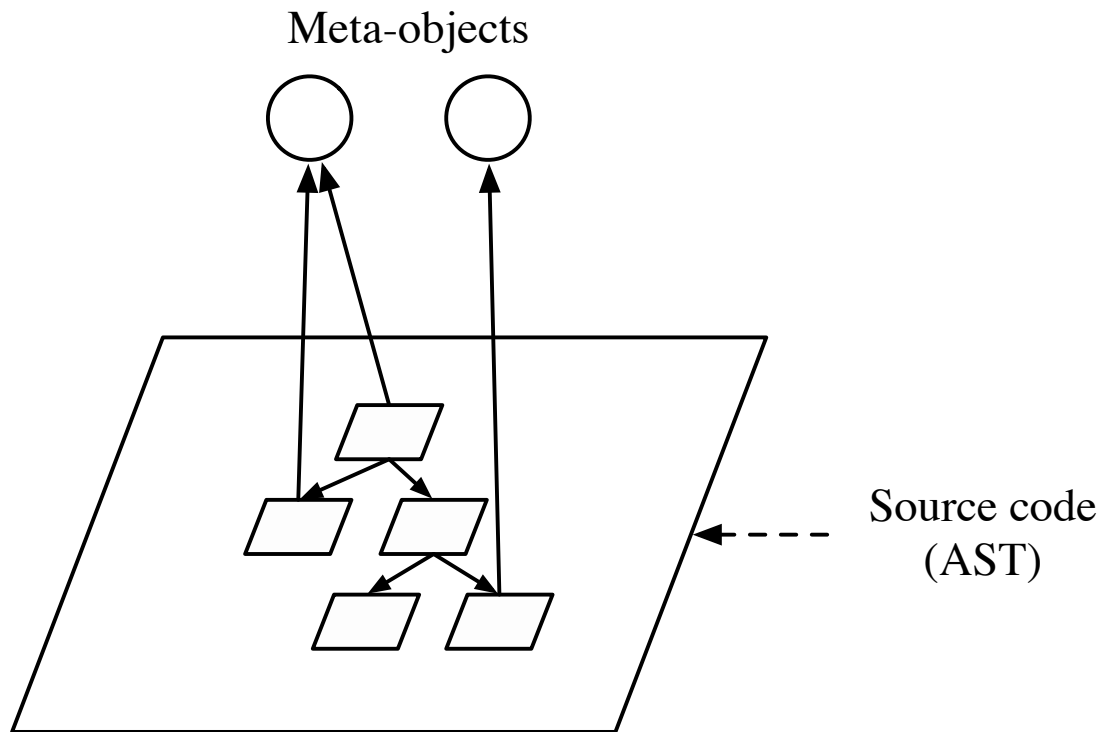
Reflectivity

Denker 2008

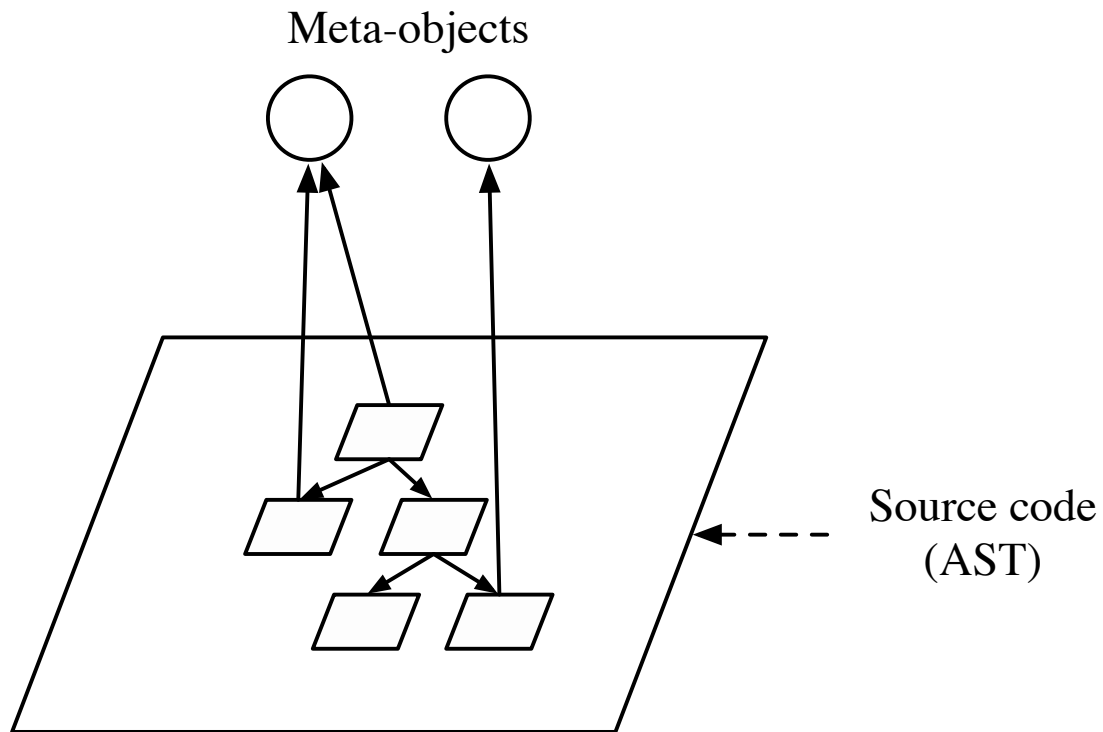
Albedo



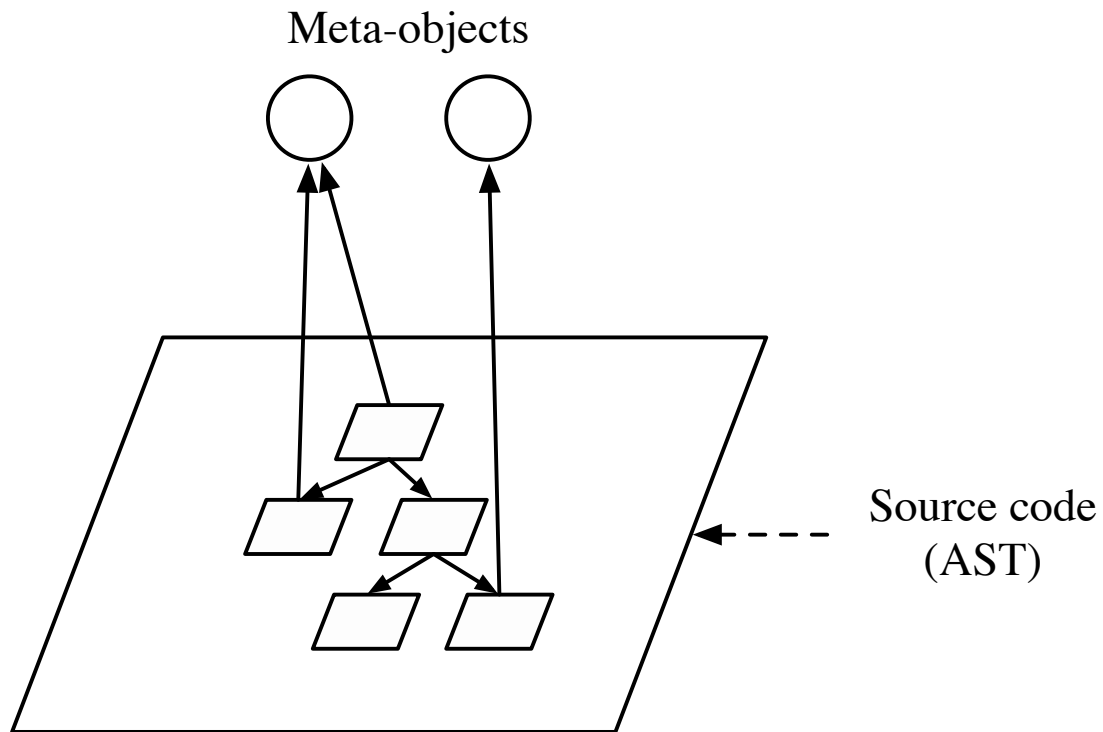
Albedo



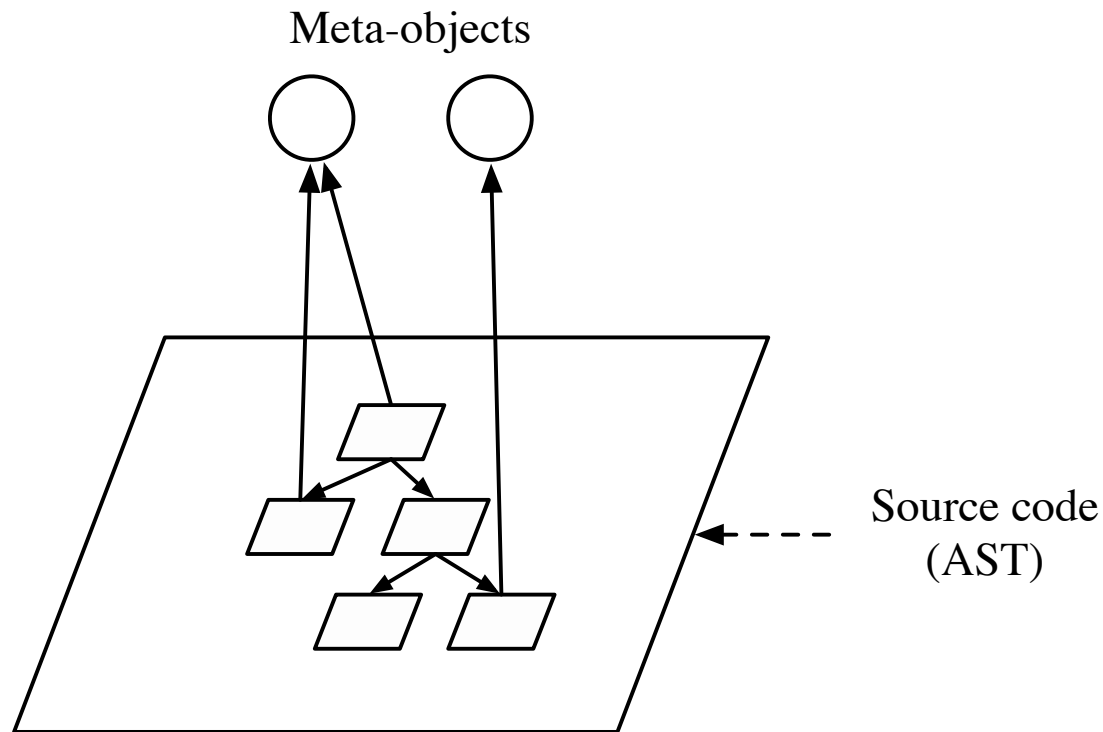
Albedo



Albedo



Albedo



Albedo
Ressia 2010

Opal Compiler

<http://scg.unibe.ch/research/OpalCompiler>